

Panel 1

Last Time:

Network Layer

Routing Shortest Path

 → Flooding

 Distance Vector

 Link State Rt.

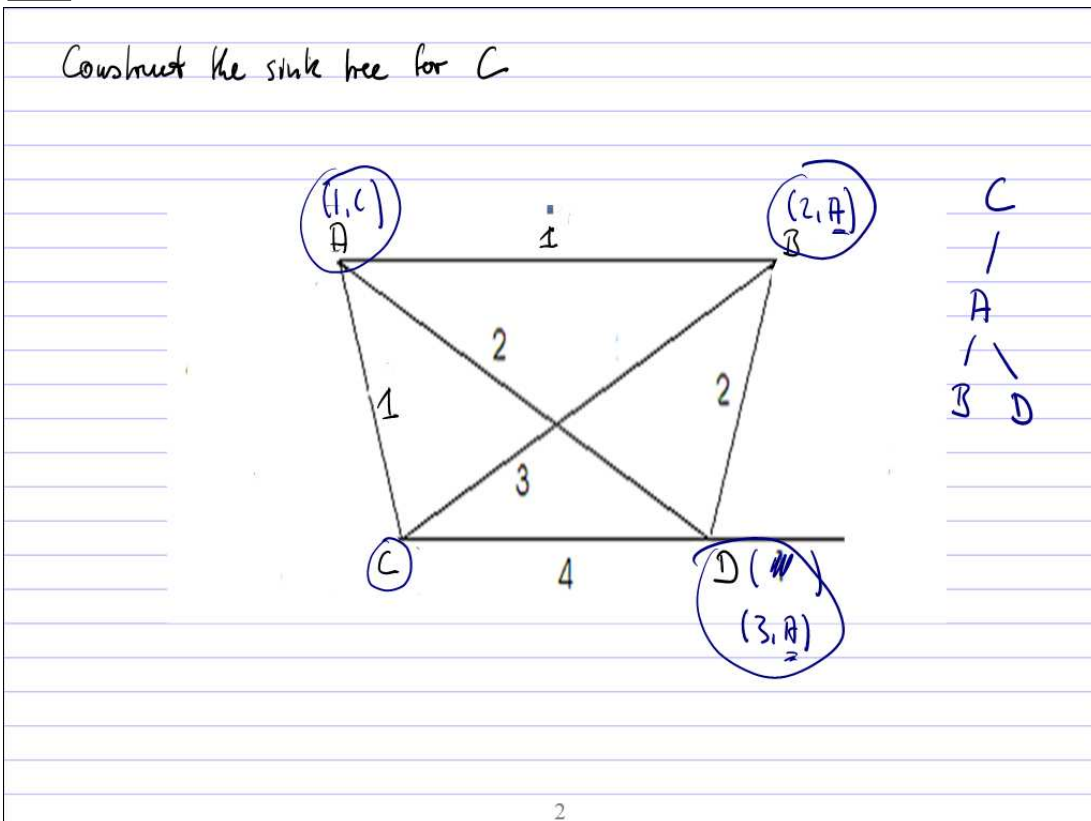
Also: Parity Bit

 Shuffling

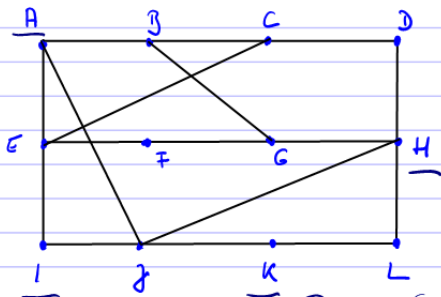
 Framing

1

Panel 2



Panel 3



Each router knows the delay to each neighbor (via ping packet).

	(H)	(H)	(I)	(K)
A	0	20	24	21
B	12	31	36	28
C	25	19	18	36
D	40	8	27	24
E	14	30	7	22
F	23	19	20	40
G	18	6	31	31
H	17	0	20	19
I	21	14	0	22
J	9	7	11	10
K	24	22	22	0
L	29	9	33	9

J		(8)
8	A	
20	A	
28	I	
20	H	
17	I	
30	I	
18	H	
12	H	12
10	I	10
0	-	
6	K	6
15	K	

Find the missing entries in the routing table for J

Panel 4

Distance vector routing problem:
 Count-to-infinity problem - explain

"Bad news travel too slow"

"Alg. converges too slowly if router / connection goes down"

Panel 5

Link State Routing

Each router must:

1. Discover its neighbors + unique network address
2. Measure delay to each neighbor
3. Construct a packet with this information
- ⊗ 4. Distribute this info to all routers
5. Compute shortest path to every router

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Panel 6

1. Discover neighbors

Protocol

Send "HELLO" packet with your address to your neighbors.

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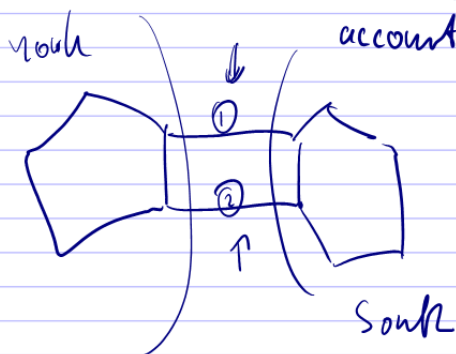
Panel 7

2. Measure Delays

Send ECHO packet, start timer.

Wait for return, divide time in half
 \approx delay

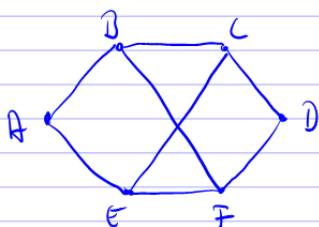
Should you take load into account or not



might get oscillation

Panel 8

3. Building Link State Packet



Distrib. these packets

- when state changes
- every 10 minutes

Panel 9

4. Distribution of Link State Packets

Tricky: Routers getting first packets will change their routes \Rightarrow different routers will have different info of network

Use flooding,
use sequence / age field to
control the flood

9

Panel 10

5. Computing the new Routes

Use Dijkstra's Algorithm

Works well, is use today (with modifications)

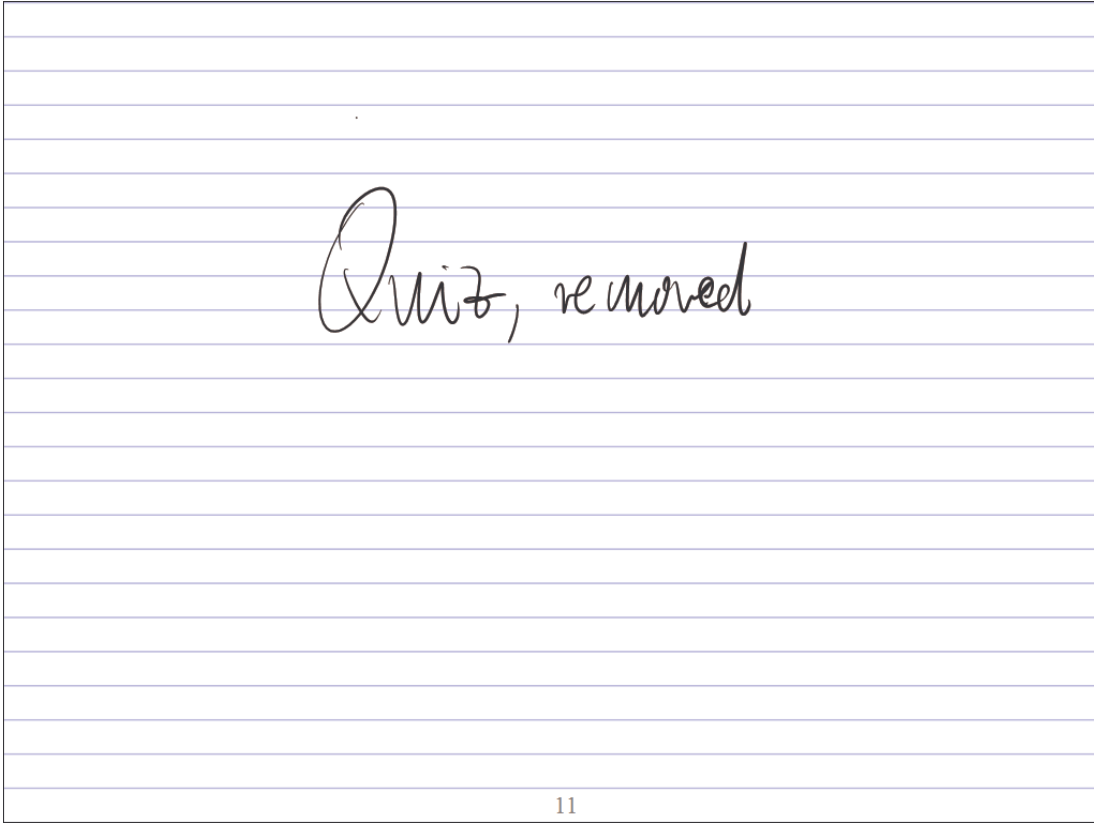
Problem: large Routing tables

\Rightarrow Use hierarchy

Thm: Use $\ln(N)$ - levels in hierarchy
for N routers!

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Panel 11



11